

A Variant of the Kochen-Specker Theorem Locating Value Indefiniteness

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Joint work with C. S. Calude and K. Svozil
arXiv:1503.01985

Between Kochen-Specker and Quantum Indeterminism

Kochen-Specker Theorem

- ▶ Impossibility of a (**complete**) noncontextual hidden variable theory for states in $d \geq 3$ Hilbert space
- ▶ “Quantum contextuality”
- ▶ State-independent

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There is a gap between this and the common interpretation of quantum measurements

Eigenvalue-Eigenstate link

A system in a state $|\psi\rangle$ has a definite property of an observable A if **and only if** $|\psi\rangle$ is an eigenstate of A .

Can we formally close this gap and show the *extent* of value indefiniteness?

The Kochen-Specker Theorem

A *context* in \mathbb{C}^n is a set of n compatible (commuting) observables.

In $n \geq 3$ Hilbert space there is a finite set of (one-dimensional projection) observables \mathcal{O} such that

the following three are in contradiction:

1. Every observable is assigned a definite value of 0 or 1;
2. These definite values are noncontextual;
3. Exactly one observable in each context is assigned the value 1.

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there is no **value assignment function** $v : \mathcal{O} \rightarrow \{0, 1\}$ with:

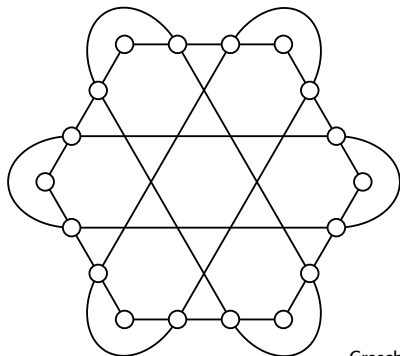
1. v is total; i.e., $v(P)$ is defined for all $P \in \mathcal{O}$;
2. v is a function of P only;
3. For every context $C \subset \mathcal{O}$:

$$\sum_{P \in C} v(P) = 1.$$

Greechie Diagrams & Value Assignments

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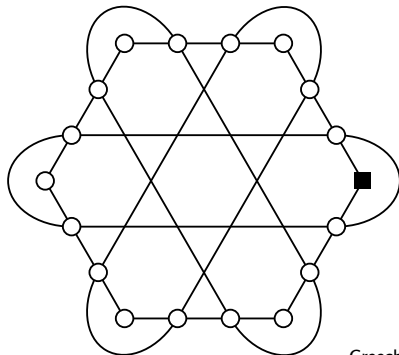
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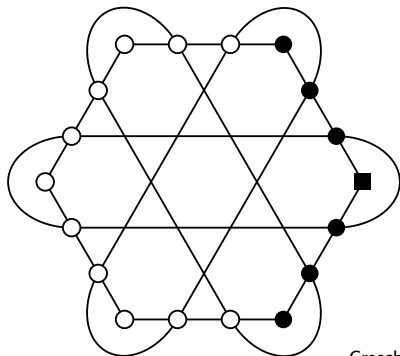
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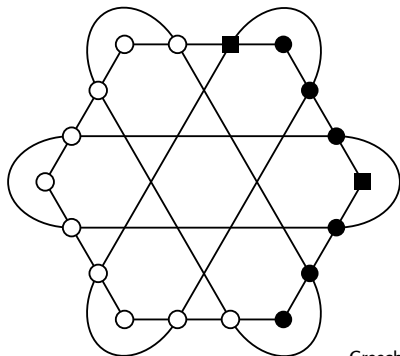
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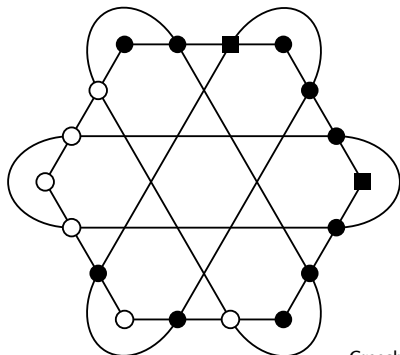
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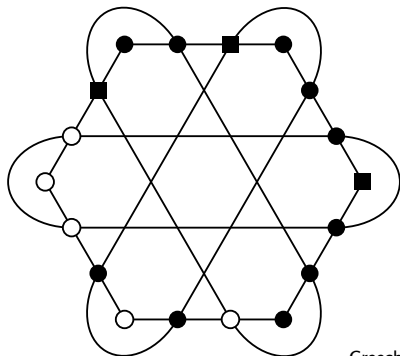
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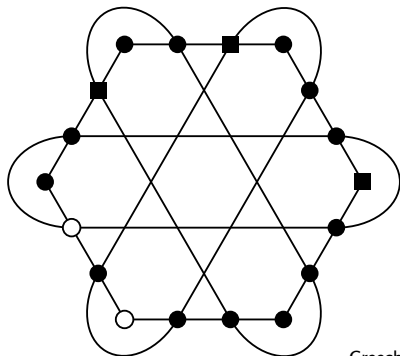
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The Extent of Value Indefiniteness

There is no value assignment function $v : \mathcal{O} \rightarrow \{0, 1\}$ with:

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Either, we reject:

- ▶ QM: But then we depart from quantum theory;
- ▶ NC: Definite values depends on measurement context;
- ▶ VD: **Some** observables are *value indefinite*.

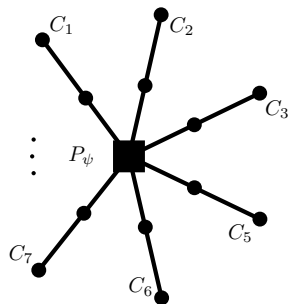
If we insist that *value definite* observables behave noncontextually, then only guaranteed that *some* observables are value indefinite

How Much Value Indefiniteness Is Reasonable?

- ▶ Can we go further and show that all observables (or how many) must be value indefinite?
- ▶ Need to localise the VD hypothesis:
 - VD: Every observable is assigned a defined value.

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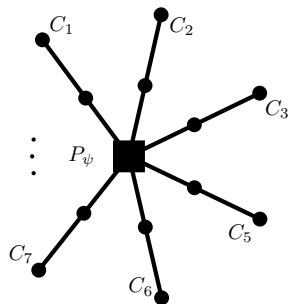
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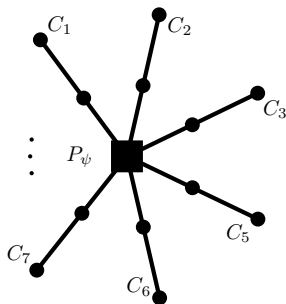
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- ▶ If system is in state $|\psi\rangle$, reasonable to expect $v(P_\psi) = 1$.
 - *One* direction of eigenvalue-eigenstate link.
- ▶ Intuitively, expect everything outside this 'star' to be value indefinite.



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 - VD: Every observable is assigned a defined value.
 - VD': **One** observable is assigned a defined value.
 - VD'': An observable P assigned 1, and a *non-compatible* observable P' value definite.
- ▶ If system is in state $|\psi\rangle$, reasonable to expect $v(P_\psi) = 1$.
 - *One* direction of eigenvalue-eigenstate link.
- ▶ Intuitively, expect everything outside this 'star' to be value indefinite.
- ▶ Need to localise all the assumptions if we wish to go further formally.



Generalising the Formal Framework

- ▶ Consider a value assignment function $v : \mathcal{O} \rightarrow \{0, 1\}$ as a representation of the system, rather than a HVT
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Admissibility of v

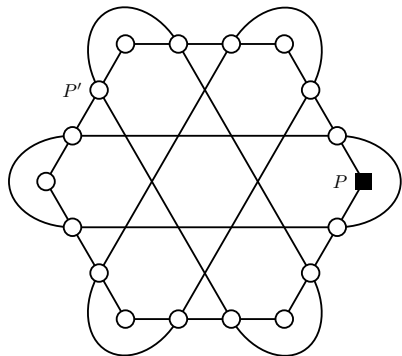
A value assignment function v is admissible if for every context $C \subset \mathcal{O}$:

- (a) if there exists a $P \in C$ with $v(P) = 1$, then $v(P') = 0$ for all $P' \in C \setminus \{P\}$;
- (b) if there exists a $P \in C$ with $v(P') = 0$ for all $P' \in C \setminus \{P\}$, then $v(P) = 1$.

Failure of Existing Greechie Diagrams

Admissibility provides a way of deducing the value definiteness of observables.

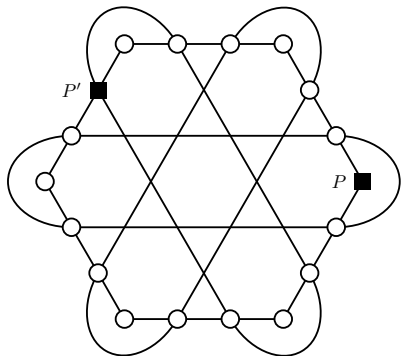
Does there exist a set of observables \mathcal{O} such that there is no admissible value assignment function with two non-compatible observables $P, P' \in \mathcal{O}$ and $v(P) = 1$ and P' value definite?



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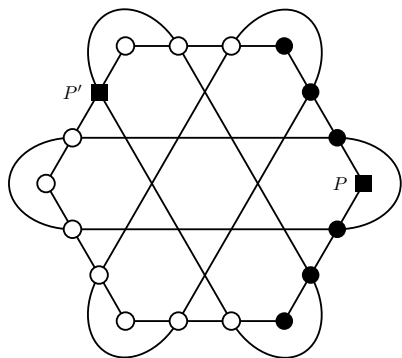
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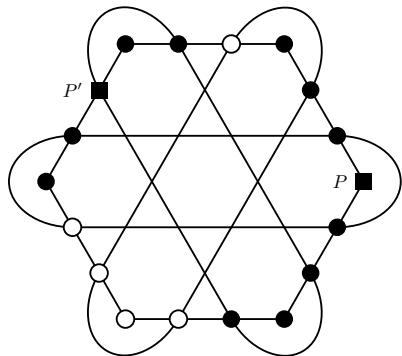
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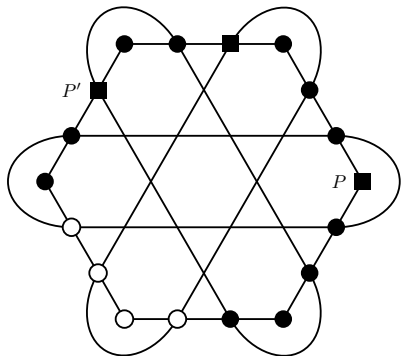
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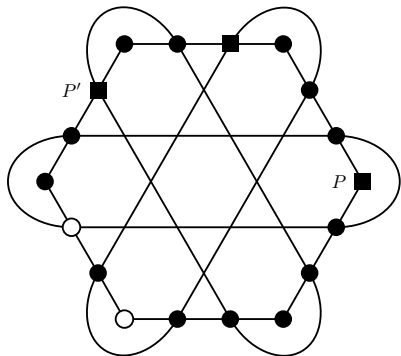
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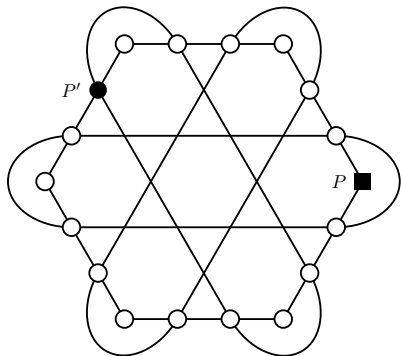
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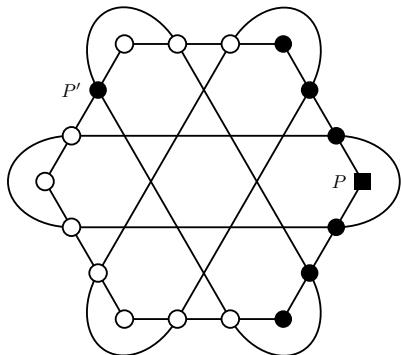


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- ▶ Need more careful interlinking of observables to obtain a contradiction.
- ▶ To prove for all P' not compatible with P we either
 - (a) Need to consider \mathcal{O} as the set of all projectors on \mathbb{C}^n ;
 - (b) Give a procedure to find $\mathcal{O} = \mathcal{O}(P')$ for a given P' .



Localised Value Indefiniteness: A Theorem

Theorem

Let $n \geq 3$. If an observable P on \mathbb{C}^n is assigned the value 1, then no other incompatible observable can be consistently assigned a definite value at all – i.e., is value indefinite.

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Let $n \geq 3$ and $|\psi\rangle, |\phi\rangle \in \mathbb{C}^n$ be states such that $0 < |\langle\psi|\phi\rangle| < 1$. Then there is a finite set of observables \mathcal{O} containing P_ψ and P_ϕ for which there is no admissible value assignment function on \mathcal{O} such that $v(P_\psi) = 1$ and P_ϕ is value definite.

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We prove in 3 steps:

1. We first prove the explicit case that $|\langle\psi|\phi\rangle| = \frac{1}{\sqrt{2}}$.
2. We prove a reduction for $0 < |\langle\psi|\phi\rangle| < \frac{1}{\sqrt{2}}$ to the first case.
3. We prove a reduction for the last case of $\frac{1}{\sqrt{2}} < |\langle\psi|\phi\rangle| < 1$ case.

► Proof

Interpretation and Conclusion

To interpret physically, need to connect P_ψ to the state of a system

Eigenstate value definiteness

If a system S is in a state $|\psi\rangle$, then $v_S(P_\psi) = 1$ for the value assignment function v_S modelling the system.

Can then conclude:

- ▶ For any state $|\psi\rangle$ can **locate** precisely the value indefinite observables
- ▶ Almost all observables are value indefinite (measure-one set)

General Interpretation

If a system is in a state $|\psi\rangle$, then the result of measuring an observable A is indeterministic unless $|\psi\rangle$ is an eigenstate of A .

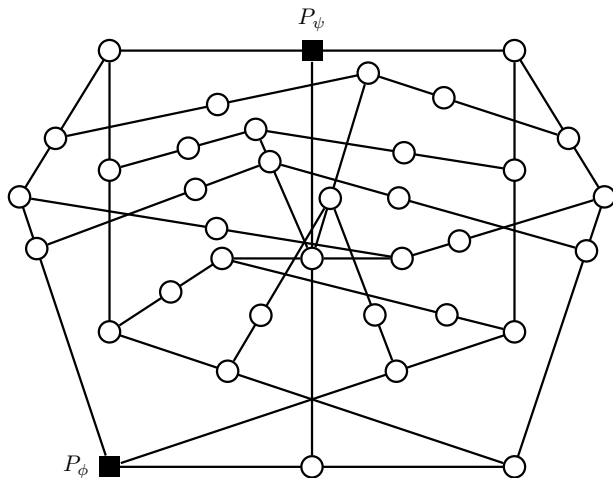
Thank You!

- ▶ A. A. Abbott, C. S. Calude & K. Svozil. *A variant of the Kochen-Specker theorem localising value indefiniteness*, [arXiv:1503.01985](#), 2015.
- ▶ S. Kochen & E. Specker. *The problem of hidden variables in quantum mechanics*, *Journal of Mathematics and Mechanics*, 17:59–87, 1967 .
- ▶ A. Cabello et al. *Bell-Kochen-Specker Theorem: A proof with 18 vectors*, *Phys. Lett. A*, 212:183–187, 1996.
- ▶ A. A. Abbott, C. S. Calude, J. Conder & K. Svozil. *Strong Kochen-Specker theorem and incomputability of quantum randomness*. *PRA*, 86(062109), 2012

Explicit Case ([▶ Back](#))

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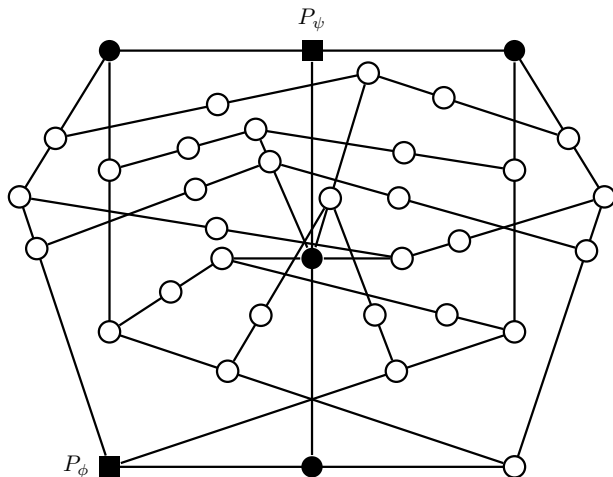
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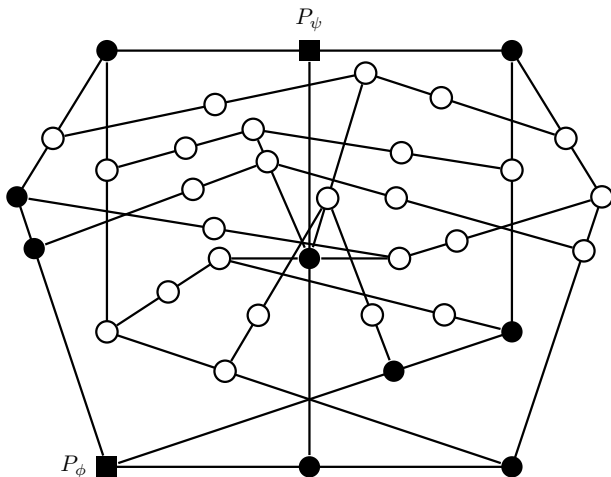
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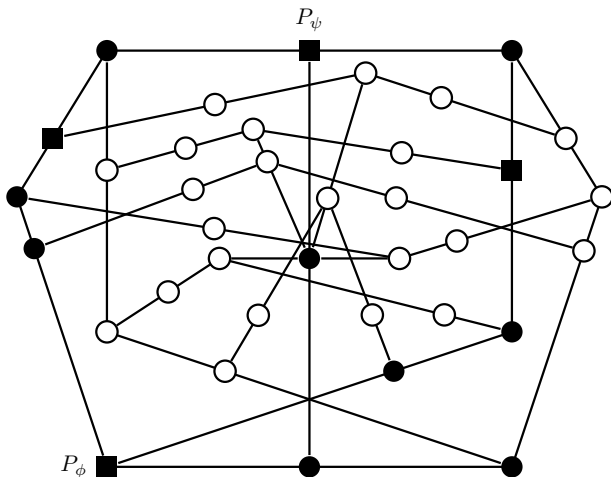
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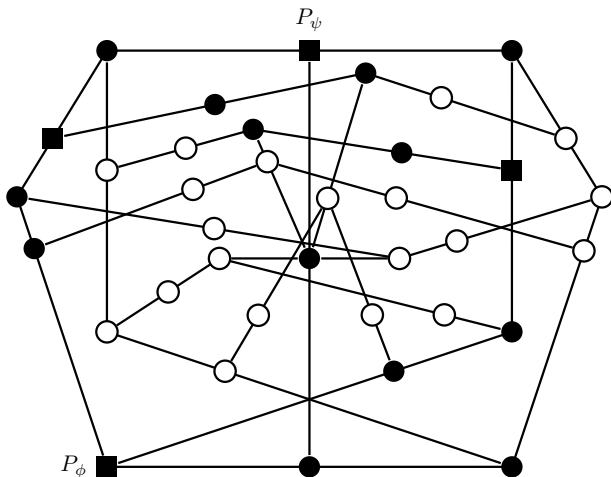
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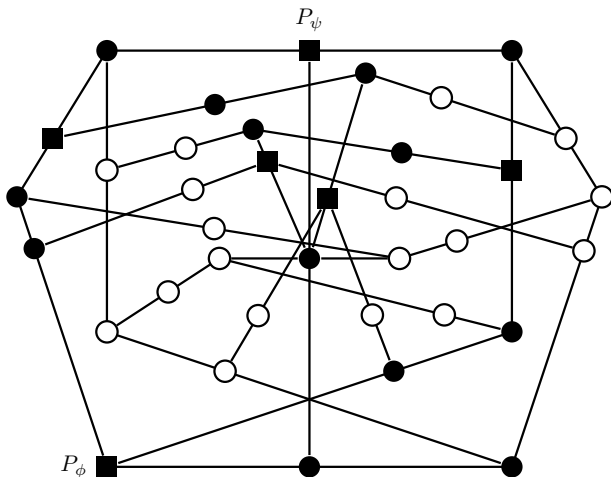
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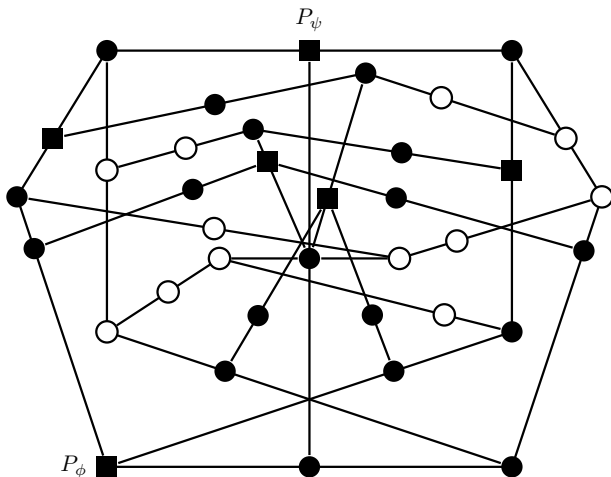
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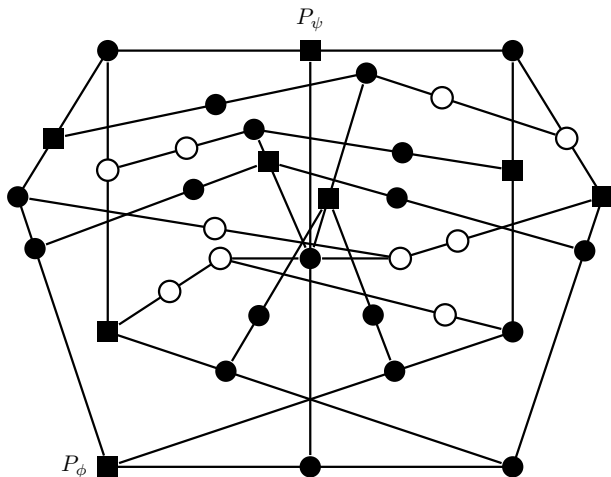
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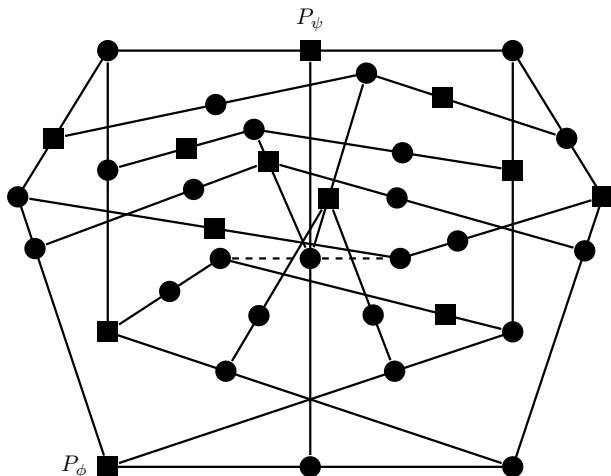
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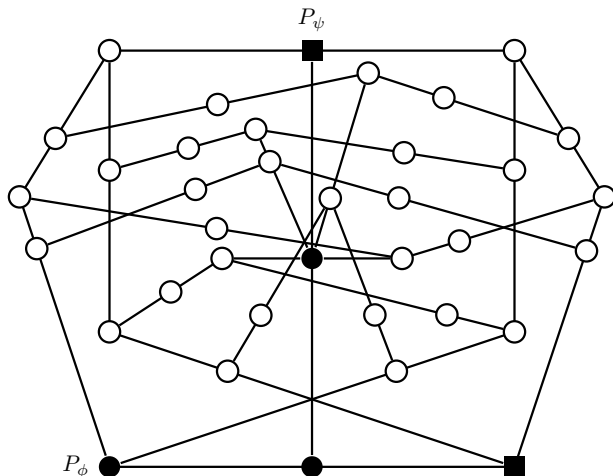
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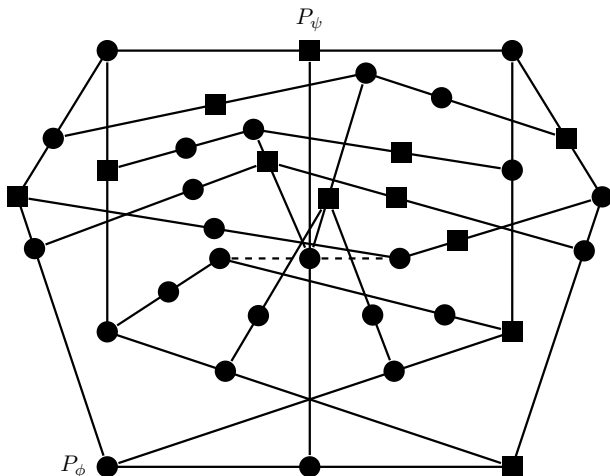
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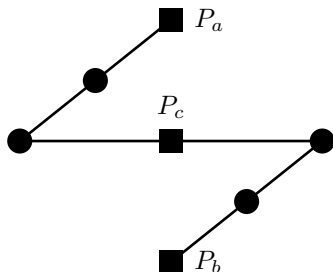
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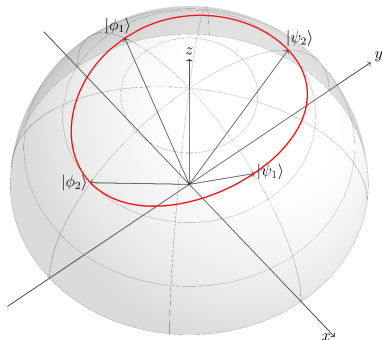
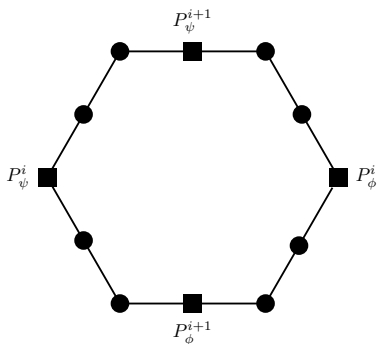
If $v(P_\psi) = v(P_\phi) = 1$ and $0 < |\langle \psi | \phi \rangle| < \frac{1}{\sqrt{2}}$, then we can find a $|\phi'\rangle$ with $\langle \psi | \phi'\rangle = \frac{1}{\sqrt{2}}$ and $v(P_{\phi'}) = 1$ under any admissible v .



The above diagram is realisable for $|\langle a|b\rangle| < |\langle a|c\rangle| < 1$.

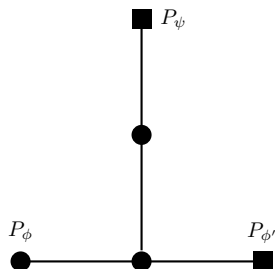
Second Reduction: Expansion [▶ Back](#)

If $v(P_\psi) = v(P_\phi) = 1$ and $\frac{1}{\sqrt{2}} < |\langle \psi | \phi \rangle| < 1$, then we can find a finite sequence of states $|\psi_1\rangle, |\phi_1\rangle; \dots, |\psi_n\rangle, |\phi_n\rangle$ such that for all i $v(P_\psi^i) = v(P_\phi^i) = 1$ and $\langle \psi_n | \phi_n \rangle = \frac{1}{\sqrt{2}}$ under any admissible v .



Completing the Proof ([▶ Back](#))

For the reductions we assumed $v(P_\phi) = 1$. If $v(P_\phi) = 0$, we can easily find $|\phi'\rangle$ with $v(P_{\phi'}) = 1$ and apply the reasoning above.



As a consequence, the set of value indefinite observables has measure 1: **almost all observables are value indefinite.**